

Zero-Trust vTPM for Confidential VMs

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Outline

- Introduction
- Our goal
- TPM-based attestation
- Zero-trust vTPM
 - Where should we run it?
 - vTPM authenticity
 - vTPM ephemeral state
 - Architecture
 - Demo
- Conclusions
- Questions

Introduction

- Protecting sensitive data end-to-end
 - At rest, in-transit and in-use
- Trusted Execution Environment (TEE)
- Confidential Virtual Machines
 - AMD SEV-SNP, Intel TDX, ARM CCA, etc
 - Protect the VM data in-use from unauthorized access
 - E.g. hypervisor / CSPs
 - Various use cases e.g. Cloud, IoT and multi-party computation

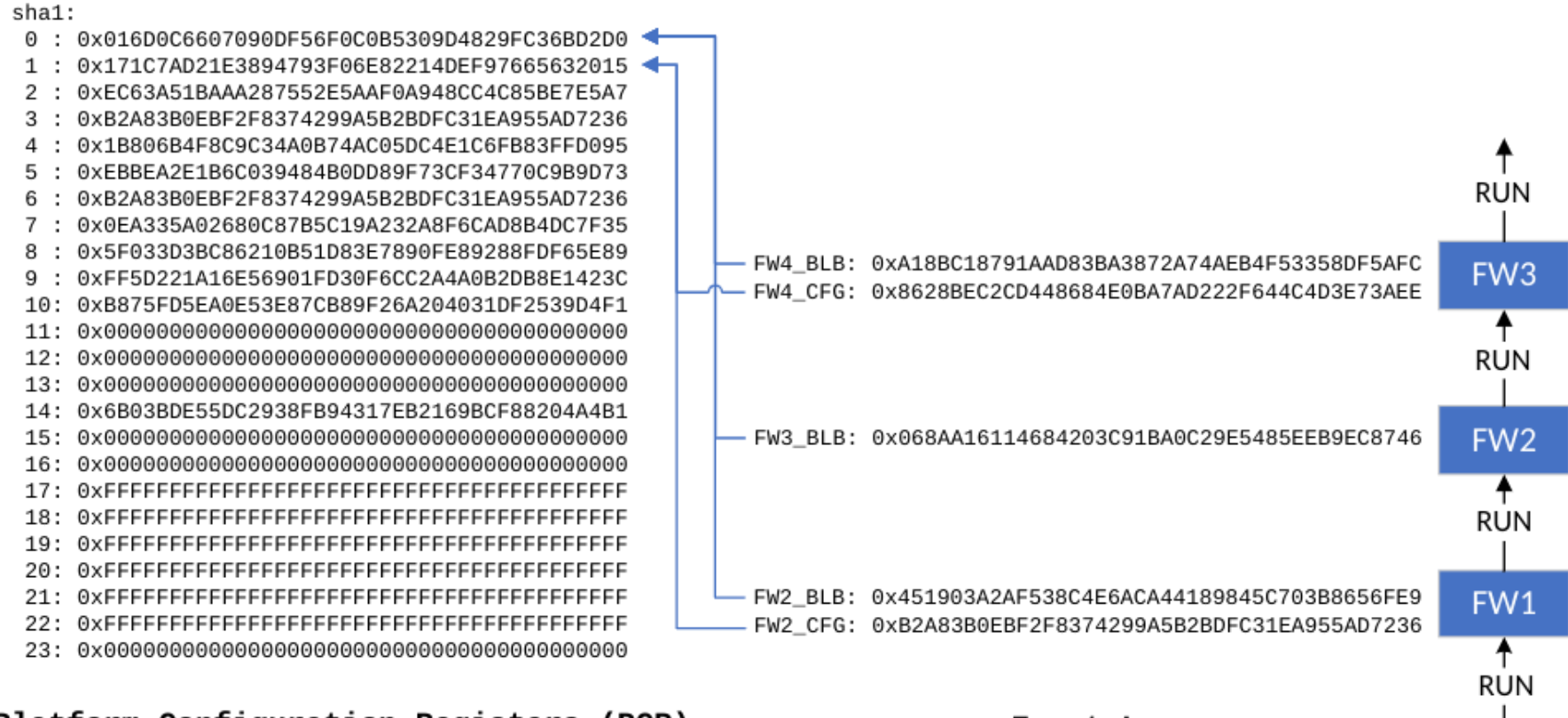
Introduction

- CVM consumption model:
 - Attest the confidential VM first
- How do we attest a CVM?
 - Boot and runtime measurements + attestation
 - VM launch measurement covers only the initial VM state
 - Solutions proposed are not standard and/or rely on untrusted entities
- Trusted Platform Module
 - Industry standard for attestation
 - Existing TPM tooling could be reused
 - Allow use of advanced attestation techniques e.g. IMA

Our Goal

- Develop a zero-trust vTPM that allow TPM-based attestation for Confidential VMs
- Target platform: AMD SEV-SNP virtual machines

TPM-based Remote Attestation



Platform Configuration Registers (PCR)

Event Log



PCR Extend operation
 $PCR_{new} := HashAlgo(PCR_{current} || digest)$

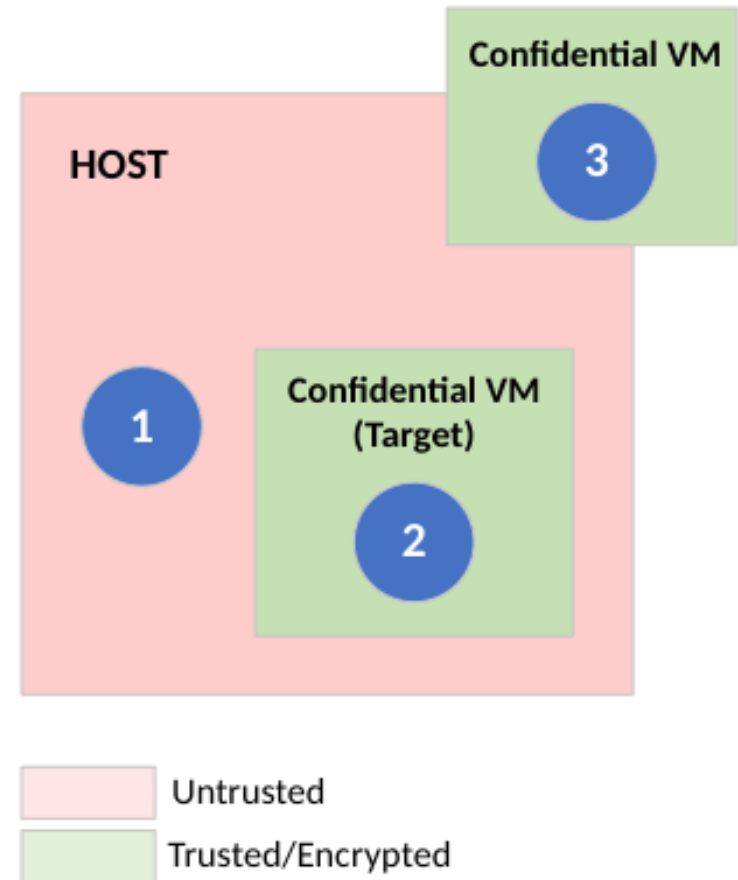
Zero-trust vTPM for attestation

1. Where should we run the vTPM?
2. How do we prove the vTPM is authentic?
3. vTPM state injection/ejection

Where should we run the vTPM?

- vTPM does not run in a physical TPM chip
 - MS TPM 2.0 Ref. implementation
 - Must protect the vTPM data from unauthorized access e.g. host OS and guest OS

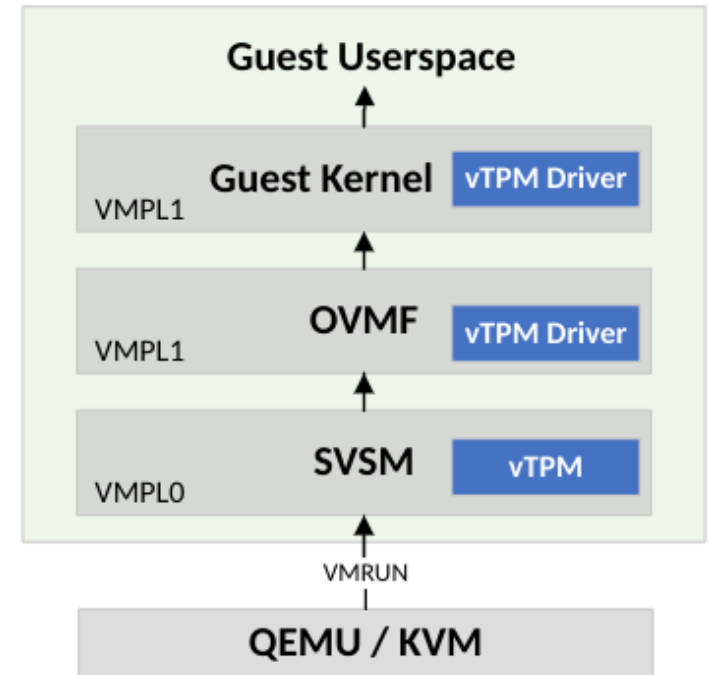
- 1 In the CVM host
- 2 In the actual CVM
- 3 In a separate CVM



Where should we run the vTPM?

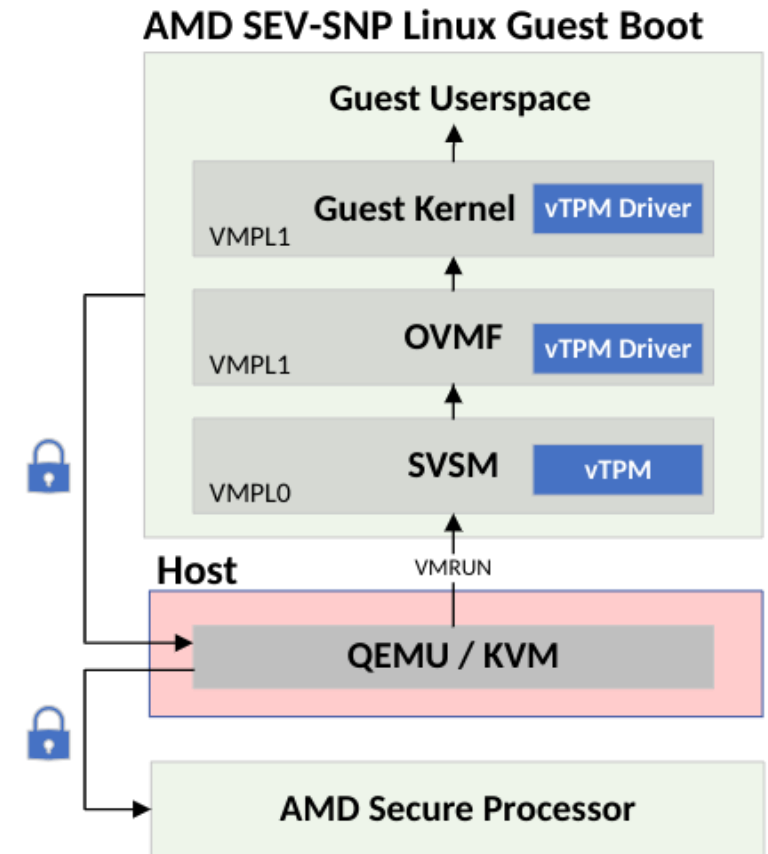
- We run the vTPM in the Secure VM Service Module
- Why do we run it in the SVSM?
 - First module that runs in SEV-SNP guests
 - Provide runtime privileged services to the guest
 - SVSM specification - <https://www.amd.com/en/developer/sev.html>
 - Leverage VM Privilege Level (VMPL) for address space isolation
- VTPM
 - Enlightened TPM drivers for SVSM
 - <https://lore.kernel.org/linux-coco/acb06bc7f329dfce21afa1b2ff080fe29b799021.camel@linux.ibm.com/>
 - [NEW] vTPM protocol – SVSM spec draft v0.62 - linux-coco mailing list
- SVSM implementations
 - <https://github.com/AMDESE/linux-svsm>
 - <https://github.com/coconut-svsm/svsm>

AMD SEV-SNP Linux Guest Boot



SVSM-vTPM Authenticity

- Physical TPM
 - Endorsement Key (EK) - TPM identity
 - EKpriv never leaves the TPM
 - EKcert: Certificate of authenticity for the EK
 - Processes used for creating and protecting the key meets the necessary security criteria (TPM Arch spec §9.5.2)
 - TPM manufacturer ships the TPM with an EK and EKcert
- SVSM-vTPM
 - MS TPM 2.0 Ref
 - Entropy source: rand() calls the rdrand assembly instruction
 - Openssl 1.1.1q
 - Create EK (TPM2_CC_CREATEPRIMARY)
 - Call the AMD-Secure Processor(SP) to issue an certificate of authenticity for the EK
 - SNP_ATTESTATION_REPORT: AMD-SP signs the attestation report with a key that chains back to AMD root key
 - Bind EK to AMD: Provide a SHA512_hash(EKpub) in the report data
 - VMPL0 Attestation report can be requested only by the SVSM (VMPCK0 key access)
 - Hypervisor cannot decrypt SNP_ATTESTATION_REPORT requests (SEV-SNP ABI spec)
 - Authentic SVSM-vTPM

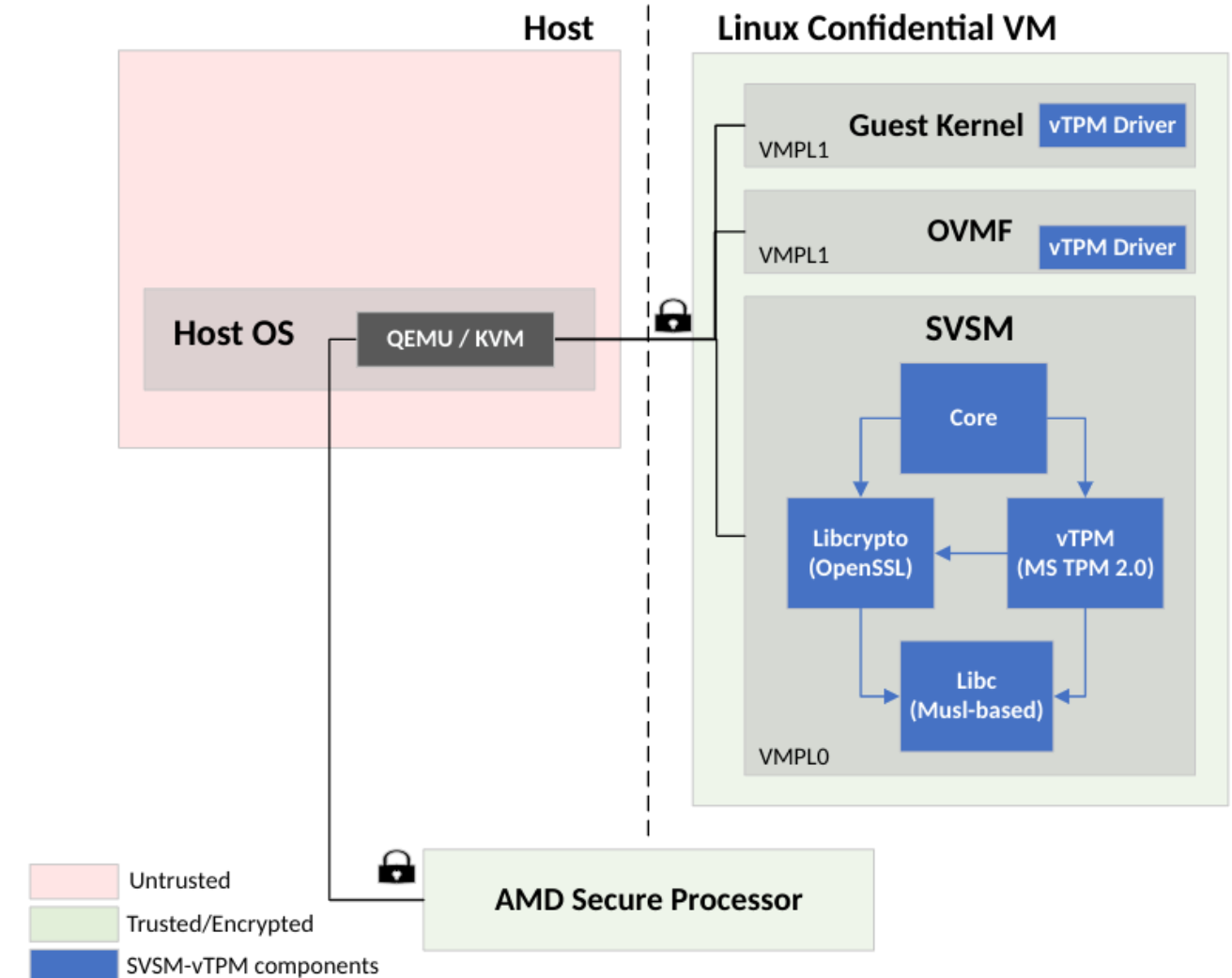


SVSM-vTPM Ephemeral State

- On every boot, we:
 - Create the EK
 - Request VMPL0 report to bind the EK to the AMD SNP platform
 - Save the VMPL0 report in the TPM NVRAM
- Advantages
 - Simple design
 - Allow TPM-based boot and runtime attestation
 - No need to eject or inject state
 - No need to early attest the VM Launch Measurement before injecting the TPM state
 - No need to secure the TPM state at rest and in-transit
 - State injection requires more code and orchestration at early boot
 - Its attack surface is considerably smaller than persistent state
 - Simplified SVSM-vTPM migration
- Disadvantages
 - May have limited use cases
 - However, full disk encryption can be enabled by an intermediate storage key

SVSM-vTPM Architecture

- SVSM-vTPM proof-of-concept
 - <https://github.com/svsm-vtpm/linux-svsm>
 - Tested the SVSM-vTPM with keylime
- Contributing it to SVSM open source projects



SVSM-vTPM Demo

- TPM NVRAM: VMPL0 Attestation Report (signed with AMD VCEK)
 - report_data: SHA512_hash(ek-pub)
 - Vmpl: 0
- Validate SVSM-vTPM
 - Attestation report signature
 - Read VMPL0 Attestation Report saved in the TPM
 - Download the VCEK certificate from the AMD website
 - Validate the attestation report signature
 - EK pub
 - Read EKpub from the SVSM-vTPM
 - Check report_data == SHA512_hash(EKpub)
 - Check VMPL == 0

Conclusions

- vTPMs can be used in CVMs to extend the chain of trust up to the guest OS
- SVSM-vTPM:
 - Protected from the guest OS and Host OS
 - Its EK uniquely identify the CVM
 - Authentic: its EKpub is chained back to AMD hardware root-of-trust
 - It's state is ephemeral (simple design) and it allows use of the TPM-based attestation
 - With small TPM driver changes the existing TPM tooling can be reused
- We're contributing the SVSM-vTPM proof-of-concept to SVSM open source projects

Questions?

- How can we improve isolation between modules in the SVSM?

Thank you!