

eBPF-based Extensible Paravirtualization



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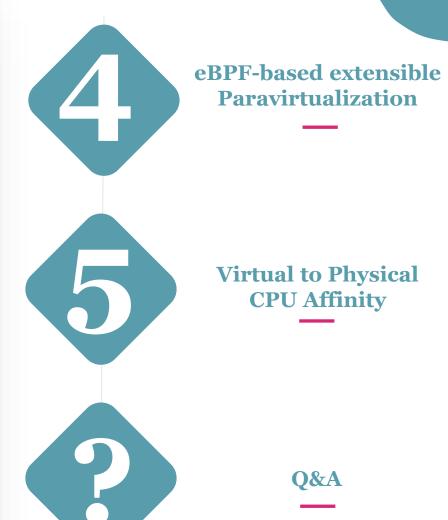


2

eBPF



Hyperup Calls



Paravirtualization

Possible Approaches



Hypercalls

Modify the guest kernel to give the hypervisor some hints



VM Introspection

The hypervisor analyzes the guest's memory



Hyperup calls



eBPF





Useful debugging tool



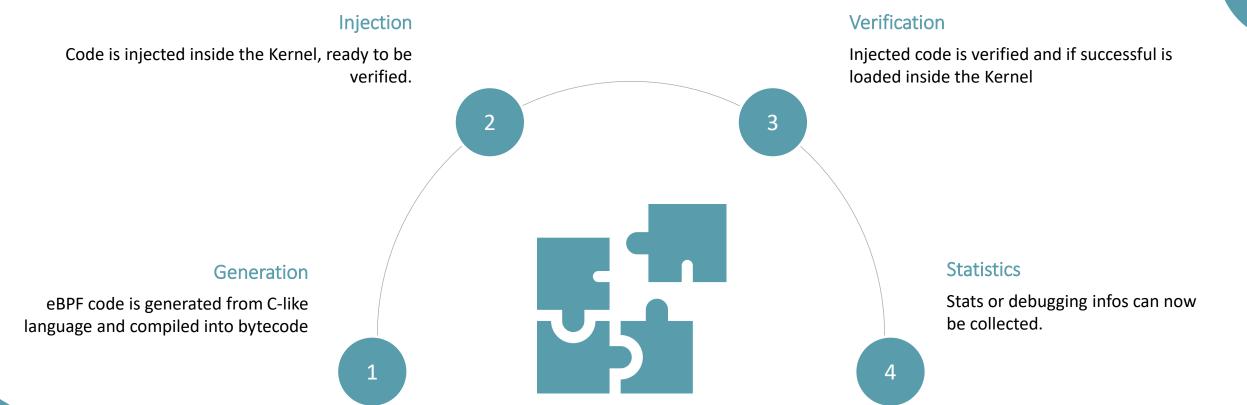
In Kernel Verifier



Kernel is not modified

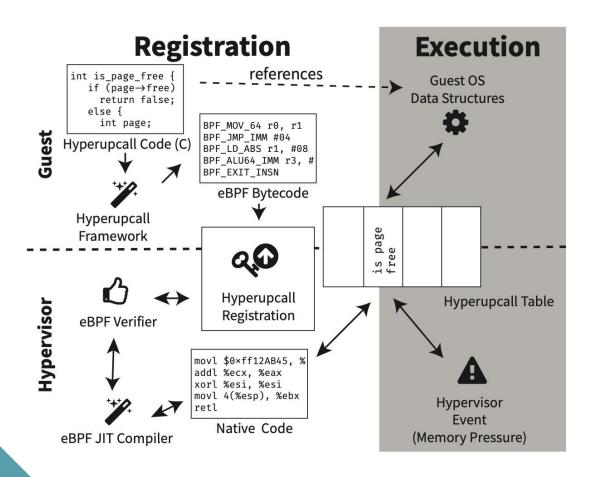


eBPF





Hyperup Calls





Performed by the Hypervisor with low overhead



Code is verified before execution



Ineffective when guest memory is encrypted



eBPF-based Extensible Paravirtualization



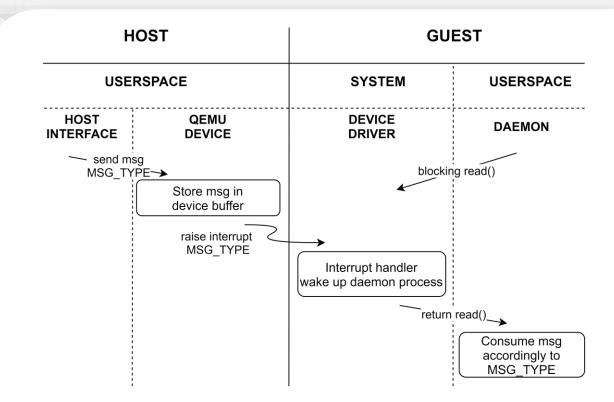
Host sends a message to the guest



Guest Agent consumes the message

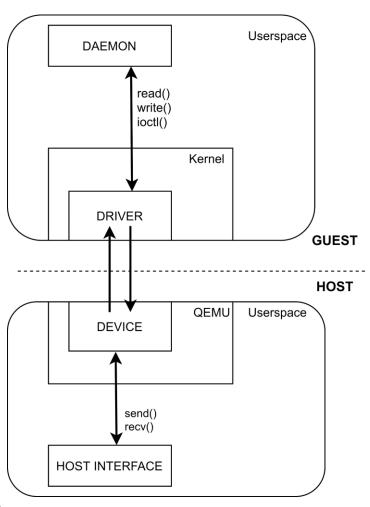


eBPF code inside the message





eBPF-based Extensible Paravirtualization



Can be loaded or unloaded at any time



No need to modify the guest's Kernel



Guest is free to decide to load the eBPF program or not.





Comparison w.r.t. Hyperupcalls

eBPF-based	Extensible	Paravirtua ¹	lization
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Host to Guest

eBPF

Async. Response

Hyperupcalls

Guest to Host

eBPF

Invoked by the HV



Virtual to Physical CPUs Affinity



Reasons



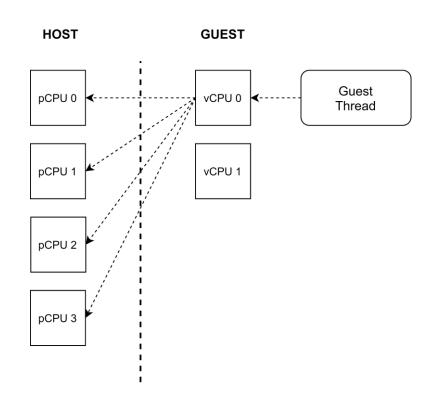
Speed

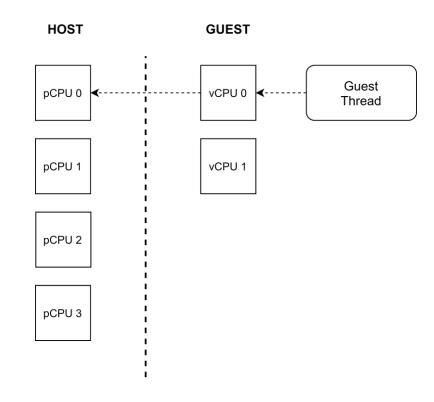




Security

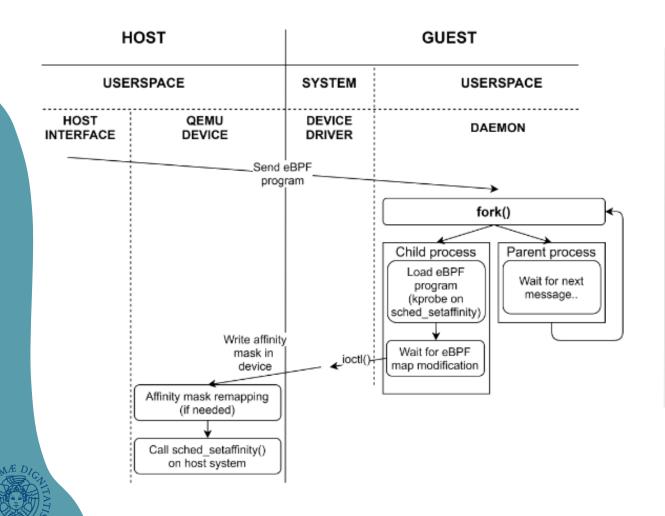
Virtual to Physical CPUs Affinity







Virtual to Physical CPUs Affinity





Kprobe on *sched_settaffinity*



Guest Agent checks for changes in eBPF map

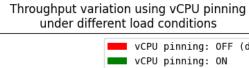


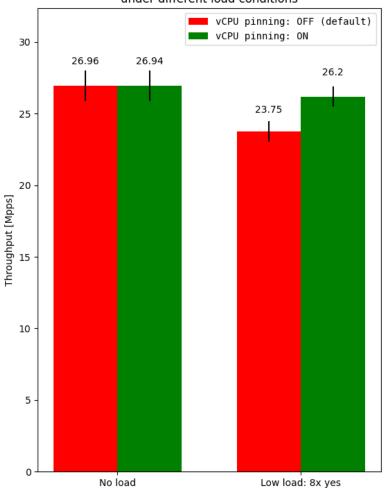
Information on bindings is sent to the hypervisor

Performance Evaluation



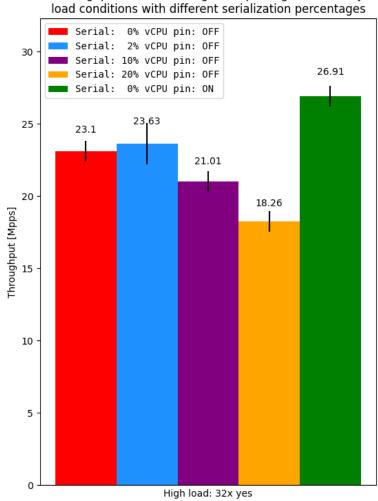
vCPU Pinning w/o HT





No Load **Low Load**

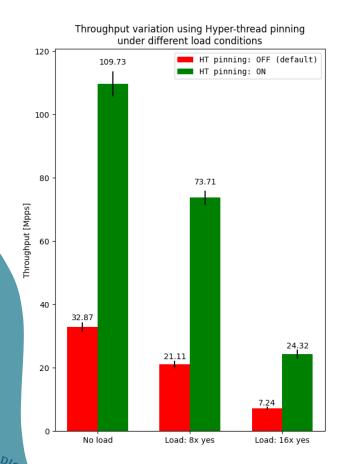


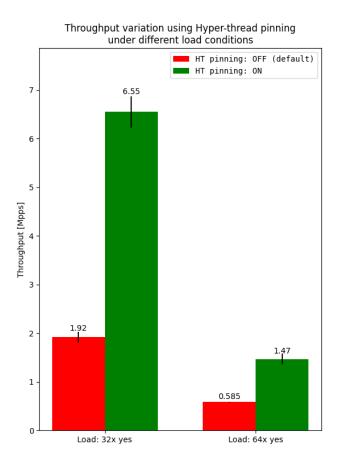


Throughput variation using vCPU pinning under heavy

Performance Evaluation

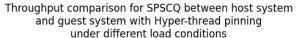
vCPU Pinning with HT

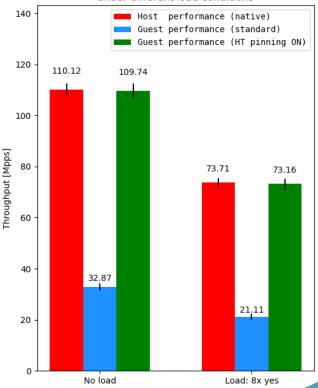






Native Perfomance







Thank You!

Any Questions?







