

Fuzzing Host-to-Guest Attack Surface in Android Protected KVM



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Agenda

- Introduction into Android pKVM
 - $\circ~$ Overview & motivation
 - Attack surface
- Virtio driver stack in pKVM
 - \circ Why fuzzing?
 - Challenges with fuzzing virtio front-end drivers
- Fuzzing virtio front-end drivers with LKL
 - Linux Kernel Library for fuzzing
 - $\circ~$ Overview of the developed fuzzers for pVM
- Conclusion & Future work

Terminology

- ABL -- Android bootloader
- AVB -- Android verified boot
- AVF -- Android virtualization framework
- GKI -- Generic kernel image
- LKL -- Linux kernel library
- Microdroid -- a Google-provided mini-Android OS that runs in a pVM
- pKVM -- Protected KVM
- pVM -- Protected virtual machine
- PVMFW -- Protected virtual machine firmware
- SMP -- Symmetric multiprocessing

Who we are

Will Deacon

- Active upstream kernel developer, co-maintaining aarch64 architecture port, locking, atomics, memory model, TLB, SMMU, ...
- Leading the Protected KVM project to enable KVM on Android



Eugene Rodionov

Android Red Team security engineer Focused on finding & exploiting vulnerabilities in low-level software in AOSP and Pixel devices



Android Protected KVM



Android Protected KVM: Overview

 Protected KVM introduces a new security model where the host and the deprivileged guest VMs mutually distrust each other.

• Mutual distrust:

- Protected KVM provides security for the guest VMs even if the host kernel is compromised
- A malicious guest cannot escape into the host (Android) or cannot compromise another guest VM \bigcirc

Deprivileged guests:

- Guest VMs don't need TrustZone privileges and run in non-secure world EL1/EL0
- **Protected KVM on Arm64: A Technical Deep Dive** by Quentin Perret
- Now You See Me, Now You Don't: Splitting pKVM Into Discrete, Mutually **Exclusive Address Spaces** by Marc Zyngier
- All Bark and no Bite: vCPU Stall Detection for KVM Guests by Sebastian Ene
- Panel discussion: KVM-based virtualization contributor Q&A by Will Deacon, et al

Building pillars of pKVM security

Host & guest VM software

- VMM and protected VM payload Ο
- Process untrusted input Ο received from the host/guest respectively
- **Prioritizing host-to-guest** Ο attacks

Hypervisor

- Enforces isolation of the guests Ο between each other and from the host
- Protects pVM bootloader and Ο sealing keys.



Attestation & Sealing

- Enable external services to attest the Ο integrity of protected VMs
- Enable per-VM instance secret data Ο



Secure boot & AVB

- Enforces authenticity of the 0 hypervisor and Android kernel
- Provides attestation services and Ο protects sealing keys

Virtio attack surface



Android Protected KVM Attack Surface



Protected VMs Non-Protected VMs Microdroid manager & payload 1 Application 0 (VM) **Guest kernel (GKI)** Microdroid bootloader (u-boot) Arbitrary kernel pvmfw (u-boot)

Guest hypercalls

Host: virtio back-end

Host <==> guest shared memory

Virtio descriptors			
buffer PA	buffer len	flags	next
0x7fff800	0x100	0×0	
• • •	• • •	• • •	
	• • •	•••	•

Virtio used buffer	Virtio avail
flags	flags
idx	idx
(desc idx, length)	desc ic
	• • •



Guest: virtio front-end

Host: virtio back-end

Host <==> guest shared memory

Virtio descriptors			
buffer PA	buffer len	flags	next
0x7fff800	0x100	0×0	
• • •	• • •	• • •	
•••	•••	•••	•





Guest: virtio front-end

1. Put request in the virtio queue

Host: virtio back-end

- 2. Process request from the queue
- 3. Put response in the virtio queue

Host <==> guest shared memory

	Virtio de	scriptors	
buffer PA	buffer len	flags	next
 0x7fff800	0x100	0x0	
• • •	• • •	• • •	
• • •	• • •	• • •	•





Guest: virtio front-end

1. Put request in the virtio queue

Host: virtio back-end

- 2. Process request from the queue
- 3. Put response in the virtio queue

	Virtio de	scriptors	
buffer PA	buffer len	flags	next
 0x7fff800	0x100	0x0	
• • •	• • •	• • •	
• • •	• • •	• • •	•



Guest: virtio front-end

- Put request in the 1. virtio queue
- 4. Process response from the virtio queue

Host: virtio back-end

- 2. Process request from the queue
- 3. Put response in the virtio queue



Host <==> guest shared memory

		Virtio de	scriptors	
	buffer PA	buffer len	flags	next
→	0xfffff00	0xffff	0×0	1:
	• • •	• • •	• • •	
	• • •	• • •	• • •	•





Guest: virtio front-end

- 1. Put request in the virtio queue
- 4. Process response from the virtio queue

Protected VM virtio attack surface



Virtio hardening in Linux mainline & u-boot

Host-to-guest attack vector **isn't new** for Linux mainline.¹

However, this attack vector **is new for Android** and pKVM, in particular. Virtio implementation in u-boot wasn't hardened against malicious host.²

Fuzzing virtio front-end drivers in the Linux kernel

Why fuzzing virtio drivers?

- One of the most effective ways to find stability and security issues in C/C++ code
- Fuzzing provides continuous security
- Fuzzer harness could be potentially reused across GKI/u-boot • as long as the same fuzzing engine is used
- Not too many security tools for Linux/Android kernel to choose from: • syzkaller¹ & syzbot² is a 'de-facto standard' fuzzing tools for kernel

[1] https://aithub.com/aooale/svzkaller

[2] https://svzkaller.appspot.com/upstream

Virtio fuzzing: challenges



LKL Overview

Linux kernel library (LKL)¹ builds Linux kernel as a user-space library

- Implemented as Linux arch-port
- LKL vs UML

LKL building blocks:

- Host environment API -- portability layer
- Linux kernel code
- LKL syscall API exposed to the user-space application

Run kernel code without launching a VM:

- kernel unit testing
- fuzzing!^{2,3}

[1] <u>https://github.com/lkl/linux</u>
 [2] Xu et al., Fuzzing File Systems via Two-Dimensional Input Space Exploration
 [3] <u>https://github.com/atrosinenko/kbdysch</u>



int ret = lkl_start_kernel(&lkl_host_ops, "mem=50M");

lkl_mount_fs("sysfs"); lkl_mount_fs("proc"); lkl_mount_fs("dev");

dev_t dev = makedev(MISC_MAJOR, UHID_MINOR); int mknod_result = lkl_sys_mknodat(AT_FDCWD, "/dev/uhid", S_IFCHR | S_IRUSR | S_IWUSR, dev);



int ret = lkl_start_kernel(&lkl_host_ops, "mem=50M");

lkl_mount_fs("sysfs"); lkl_mount_fs("proc"); lkl_mount_fs("dev");

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Anatomy of LKL fuzzer

LKL enables fuzzing Linux kernel code in user-space

• use in-process fuzzing engine, such as libFuzzer

Advantages:

- high fuzzing performance on x86_64 cores
- lightweight fuzzers (no need to run VMs)
- easy debugging & crash reproducing (i.e. gdb)
- hardware emulation (e.g. PCI)

Limitations:

- no SMP in LKL
- x86_64 vs aarch64 -- potential false positives, true negatives



Virtio front-end fuzzers

Kernel under test:

• android13-5.10

virtio_ring:

- fuzzes ring-buffer processing functionality
- handles both split & packed mode

virtio_pci:

- fuzzes PCI configuration space
- LKL arch-specific implementation of PCI bus
- mock-out PCI MMIO in the fuzzer harness

virtio_blk:

• mutates the virtio_blk configuration block



Virtio_blk fuzzer finding

int block_read_full_page(struct page *page, get_block_t *get_block) struct buffer_head *bh, *head, *arr[MAX_BUF_PER_PAGE];

```
• • •
do
      (buffer_uptodate(bh))
  if
     continue;
```

```
if (!buffer_mapped(bh)) {
  int err = 0;
```

```
if (buffer_uptodate(bh))
     continue;
}
```

```
OOB write on
stack
```

```
arr[nr++] = bh;
```

} while (i++, iblock++, (bh = bh->b_this_page) != head);

Virtio blk fuzzer finding

With the block size 0xe5e5e5e5:

o `inode->i_blkbits == 32`

- o `1 << READ_ONCE(inode->i_blkbits)` is undefined behavior in C
- o `1 << READ_ONCE(inode->i_blkbits) == 1` on x86 architecture

```
static struct buffer_head *create_page_buffers(struct page *page, ...)
```

BUG_ON(!PageLocked(page));

if (!page_has_buffers(page))

```
create_empty_buffers(page, 1 << READ_ONCE(inode->i_blkbits), b_state);
return page_buffers(page);
```

Fuzzing virtio driver stack in u-boot

- Both pvmfw (1st stage) and microdroid bootloader (2nd stage) are based on u-boot
 - rely on virtio blk to get boot configuration and virtio console for debug output
- Fuzzing and ASAN for sandbox patch series enable fuzzing for virtio front-end drives:
 - works for u-boot in sandbox mode
 - provide coverage-guided libFuzzer-based fuzzing \bigcirc
 - enables AddressSanitizer for the fuzz target \bigcirc

• Findings:

• virtio: Harden and test vring patch series

Fully controlled OOB write in u-boot

```
static void detach_buf(struct virtqueue *vq, unsigned int head)
  while (vq->vring.desc[i].flags & nextflag) {
    virtqueue_detach_desc(vq, i); // <== i is 00B</pre>
    i = virtio16_to_cpu(vq->vdev, vq->vring.desc[i].next);
    vq->num_free++;
  • • •
int bounce_buffer_stop(struct bounce_buffer *state)
  // state is 00B and point to a fully attacker-controlled memory
  if (state->flags & GEN_BB_WRITE)
    memcpy(state->user_buffer, state->bounce_buffer, state->len);
  free(state->bounce_buffer);
  return 0;
```



Conclusion



LKL-based virtio fuzzers continuously run in Google's internal ClusterFuzz engine.



Virtio fuzzing effort

led to identification and proactive mitigation of multiple security and stability issues in GKI and u-boot.



Need your support in **improving fuzzing** for virtualized interfaces.

Future work

• Write more fuzzers targeting virtio front-end and PCI drivers

- Upstreaming LKL to Linux mainline:
 - first attempt in 2015
 - restarted in 2020^[1] -- still ongoing to integrate LKL as a submodule of UML
- Currently focusing on upstreaming LKL to Android Common Kernel mainline:
 - effort to upstream LKL as a separate architecture is WIP
 - share LKL fuzzing work with the open-source community \bigcirc



Thank you!







Appendix



Android Protected KVM: Motivation

Non-secure world

Secure world



Prioritizing host-to-guest attacks in pKVM

Guest-to-host VM escapes is a traditional threat model for modern VMMs and hypervisors. Android Virtualization Framework in Android 13 doesn't allow running arbitrary guest VMs.

LKL KASan details

KASan provides actionable reports for invalid memory access:

- OOB, user-after-free, double-free
- covers stack, heap and globals

User-space ASan in LKL:

- ASan shadow memory poisoning routines are invoked in global constructors
- Which might be problematic due to specifics of globals initialization in Linux kernel

LKL implements generic KASan:

- -fsanitize=kernel-address
- arch-specific KASan implementation



LKL fuzzing coverage

LKL relies on libFuzzer-based fuzzing code coverage instrumentation.

KCOV is an alternative solution:

 needs additional implementation to feed the coverage feedback to
 libFuzzer engine

	Compile Line kernel sourc
-f	sanitize=fuzzer-
	Incremental I of IkI.o (vmlin
	lkl.o (vmlinu



How to develop an LKL fuzzer

• Identify an interface to fuzz

• use 'realistic' attack surface (i.e. reachable from user-space or from the hardware)

- Enable the kernel feature under test in the kernel config which doesn't depend on aarch64 features or SMP
- Mock-out low-level interfaces if needed LKL already comes with virtio back-end and arch-specific PCI implementations
- Provide fuzzer harness which sends the fuzzer's entropy to the target kernel interface

Output of virtio_blk fuzzer

./virtio_blk-fuzzer -close_fd_mask=3

cov: 3662 ft: 6239 corp: 92/178b lim: 4 exec/s: 455 rss: 96Mb L: 2/4 MS: 1 #455 NEW cov: 3662 ft: 6248 corp: 93/180b lim: 4 exec/s: 472 rss: 96Mb L: 2/4 MS: 2 #472 NEW cov: 3662 ft: 6249 corp: 94/184b lim: 4 exec/s: 495 rss: 96Mb L: 4/4 MS: 3 #495 NEW #496 NEW cov: 3662 ft: 6250 corp: 95/185b lim: 4 exec/s: 496 rss: 96Mb L: 1/4 MS: 1 cov: 3662 ft: 6252 corp: 96/188b lim: 4 exec/s: 510 rss: 96Mb L: 3/4 MS: 4 #510 NEW #511 NEW cov: 3662 ft: 6260 corp: 97/190b lim: 4 exec/s: 511 rss: 96Mb L: 2/4 MS: 1 #521 NEW cov: 3662 ft: 6261 corp: 98/194b lim: 4 exec/s: 521 rss: 96Mb L: 4/4 MS: 5 #525 NEW cov: 3662 ft: 6267 corp: 99/198b lim: 4 exec/s: 525 rss: 96Mb L: 4/4 MS: 4

• • •

• • •